

# Formal Analysis about Security Requirements of a Group Authentication Protocol by Scyther

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# Scyther

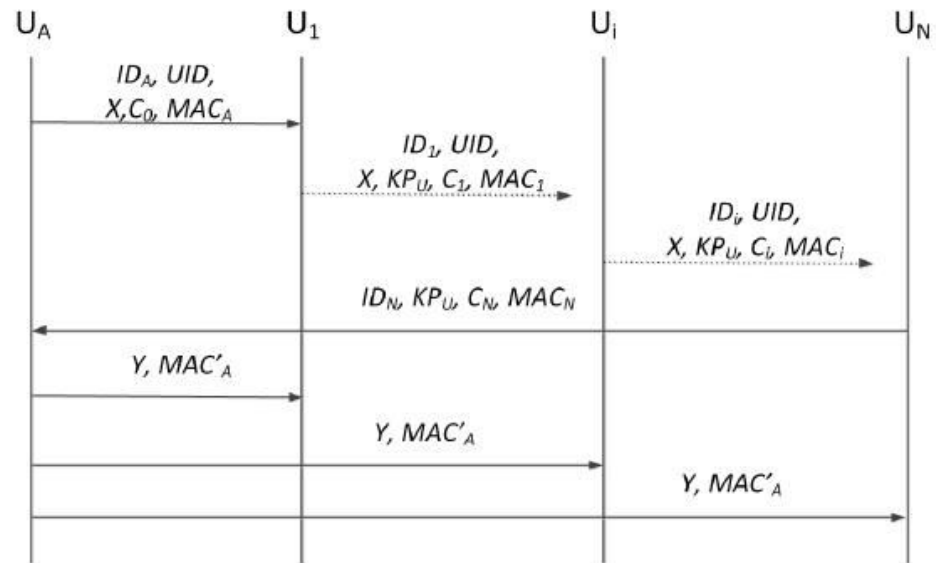
- Scyther is a tool for the automatic verification of security properties
  - <http://www.cs.ox.ac.uk/people/cas.cremers/scyther/>
- Claims:  $\text{claim}(R, \text{SKR}, \text{rt})$
- Commitment:  $\text{claim}(R_1, \text{Commit}, R_2, \text{rt})$
- Match( $p_1, p_2$ )

```
usertype SharedSecret;
hashfunction H;
protocol Example(A, B){
  role A{
    fresh x : SharedSecret;
    var y : nonce;
    send_1(A, B, {x}k(A, B));
    recv_2(B, A, {H(y)}sk(B));
  };
  role B{
    fresh y : nonce;
    var x : SharedSecret;
    recv_1(A, B, {x}k(A, B));
    send_2(B, A, {H(y)}sk(B));
  };
}
```

# Adversary Model

- Adversary model: Dolve-Yao model
- Adversary ability: eavesdrop, delete message, learn knowledge, create and insert messages

# Group Authentication Framework Introduction



1)  $U_A \rightarrow U_1 : ID_A, UID, X, C_0, MAC_A$ .

2)  $U_i \rightarrow U_{i+1} : ID_i, UID, X, KP_U, C_i, MAC_i$ , where  
 $1 \leq i \leq N - 1$ .

3)  $U_N \rightarrow U_A : ID_N, KP_U, C_N, MAC_N$ .

4)  $U_A \rightarrow U : Y, MAC'_A$ .

# Protocol Formalization (1)

- Discrete Logarithm Problem (DLP) and Elliptic Curve Discrete Logarithm Problem (ECDLP) based: Type 1 and Type 2
  - Only formalize DLP-based protocol of type 1
- Group member: 3
- Hard problems: type “hashfunction”
  - Diffie-Hellman problem, hash function, proxy encryption, MAC

## Protocol Formalization (2)

- Security requirements: claims
  - Mutual authentication, implicit authentication, against impersonation attack, against passive adversaries

1)  $match(h_i, H(U_A, U_i, x_a, t_i))$ , where  $1 \leq i \leq 3$ .

5)  $claim(U_A, Commit, U_i, V_i)$ , where  $1 \leq i \leq 3$ .

6)  $claim(U_R, SKR, K_G)$ , where  $R \in \{U_A, U_i\}$  and  $1 \leq i \leq 3$ .

7)  $claim(U_A, SKR, h(g(n_i), m_i))$ , where  $1 \leq i \leq 3$ .

8)  $claim(U_i, SKR, h(g(m_i), n_i))$ , where  $1 \leq i \leq 3$ .

9)  $claim(U_i, SKR, h(g(n_j), n_i))$ , where  $1 \leq i, j \leq 3$  and  $i \neq j$ .

# Results

Claim				Status	Comments
Group_authentication_DLP	UA	Group_authentication_DLP,UA1	SKR KG	Ok	Verified No attacks.
		Group_authentication_DLP,UA2	SKR $h(gn1,m1)$	Ok	Verified No attacks.
		Group_authentication_DLP,UA3	SKR $h(gn2,m2)$	Ok	Verified No attacks.
		Group_authentication_DLP,UA4	SKR $h(gn3,m3)$	Ok	Verified No attacks.
U1		Group_authentication_DLP,U11	SKR KG	Ok	Verified No attacks.
		Group_authentication_DLP,U12	SKR $h(gm1,n1)$	Ok	Verified No attacks.
		Group_authentication_DLP,U13	SKR $h(gn2,n1)$	Ok	Verified No attacks.
		Group_authentication_DLP,U14	SKR $h(gn3,n1)$	Ok	Verified No attacks.
U2		Group_authentication_DLP,U21	SKR KG	Ok	Verified No attacks.
		Group_authentication_DLP,U22	SKR $h(gm2,n2)$	Ok	Verified No attacks.
		Group_authentication_DLP,U23	SKR $h(gn1,n2)$	Ok	Verified No attacks.
		Group_authentication_DLP,U24	SKR $h(gn3,n2)$	Ok	Verified No attacks.
U3		Group_authentication_DLP,U31	SKR KG	Ok	Verified No attacks.
		Group_authentication_DLP,U32	SKR $h(gm3,n3)$	Ok	Verified No attacks.
		Group_authentication_DLP,U33	SKR $h(gn1,n3)$	Ok	Verified No attacks.
		Group_authentication_DLP,U34	SKR $h(gn2,n3)$	Ok	Verified No attacks.



THANK YOU!