

Ethical hacking

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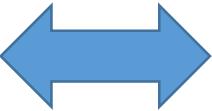
About myself

- PhD from Computer Security (Software vulnerability exploitation)
 - Research on Software vulnerabilities (error finding and exploitation)
 - Research on Sophisticated Malwares
-
- Penetration test experiences
 - Teaching Ethical Hacking (EC Council – Certified Ethical Hacker)
 - Courses on exploit writing (hardcore hacking)

Schedule

- Wednesday 10.30-12.30 Ethical hacking in general, practical tricks
- Friday 10.30-12.30 Research on memory corruption

What is ethical hacking?

- Legal (contract)
 - Promote the security by showing the vulnerabilities
 - Find all vulnerabilities
 - Without causing harm
 - Document all activities
 - Final presentation and report about the vulnerabilities
- 
- Illegal
 - Steal information, modify data (e.g. deface), make service unavailable
 - Find only the weakest link to achieve the aim
 - Do not care if the action destroys the system
 - Without documentation
 - Without report, delete all clues

Hiding during the process?

Ethical hacking concepts

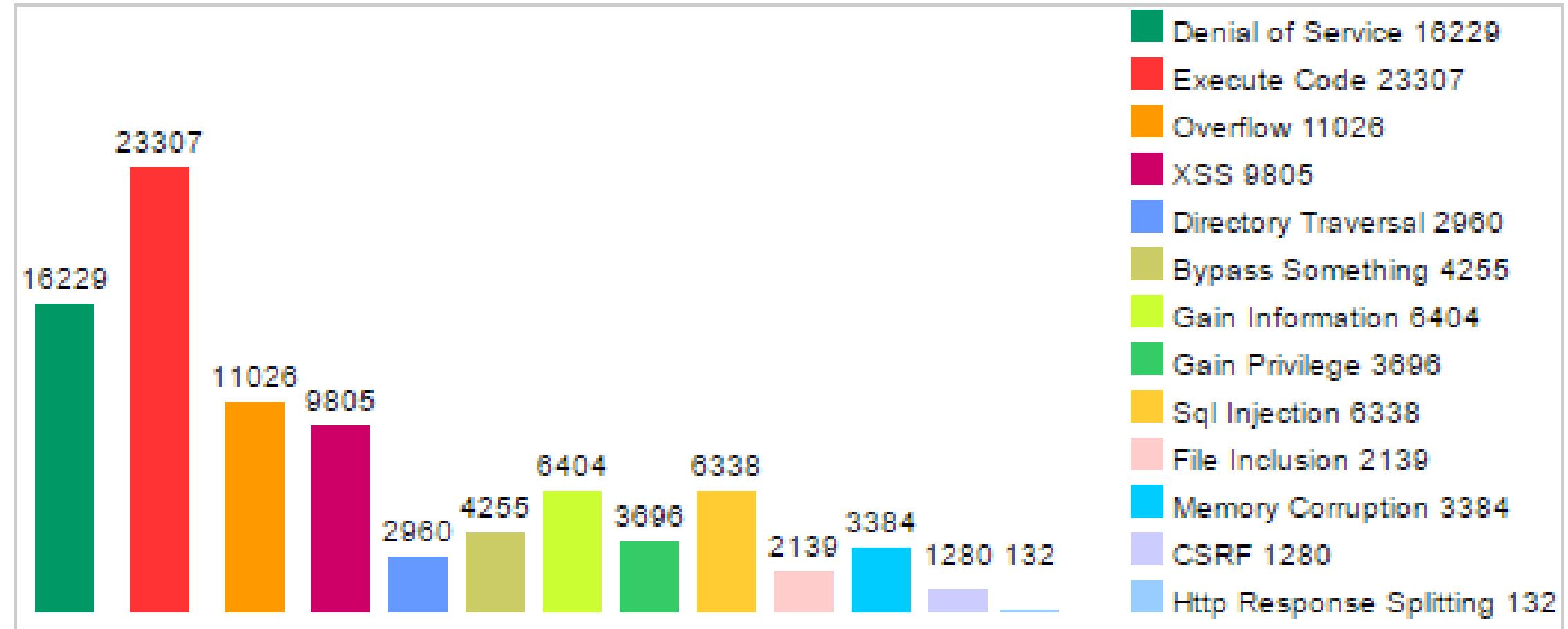
	Black box concept	Grey box concept	White box concept
Internal penetration test	X	X	X
External penetration test	X	X	X
Web hacking	X	X	X
Wireless hacking	X	X	X
Social engineering		X	

Ethical hacking steps

- General information gathering
- Technical Information gathering
- Looking for available hosts
- Looking for available services
- Manual testing
- Automatic testing
- Exploitation
- Covering tracks

Vulnerability types

Vulnerabilities By Type



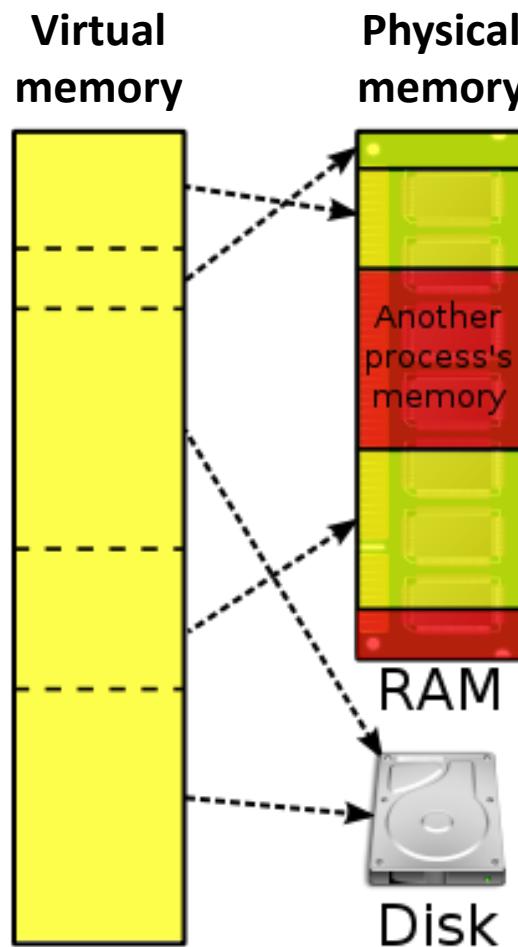
CEH topics

- Introduction to Ethical Hacking
- Footprinting and Reconnaissance
- Scanning Networks
- Enumeration
- System Hacking
- Malware Threats
- Sniffing
- Social Engineering
- Denial of Service
- Session Hijacking
- Hacking Webservers
- Hacking Web Applications
- SQL Injection
- Hacking Wireless Networks
- Hacking Mobile Platforms
- Evading IDS, Firewalls, and Honeypots
- Cloud Computing
- Cryptography

Ethical hacking course at UiA

http://ethical_hacking.project.uia.no

Virtual address space



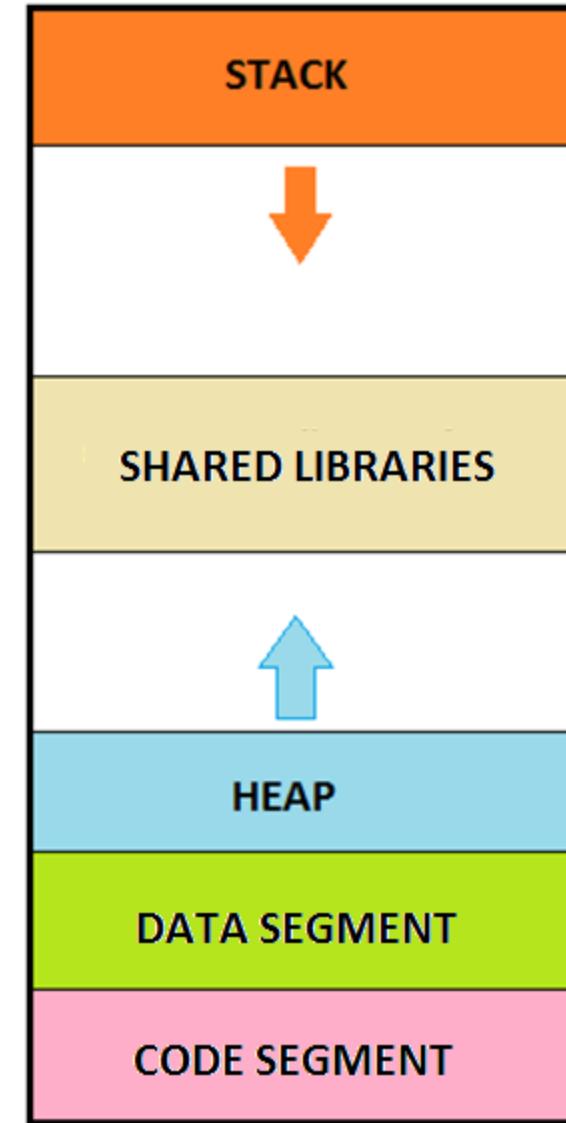
Local variables,
method
parameters,
exception handling
data,
return addresses

Dynamically
linked shared
libraries (libc)

Dynamic variables

Global variables

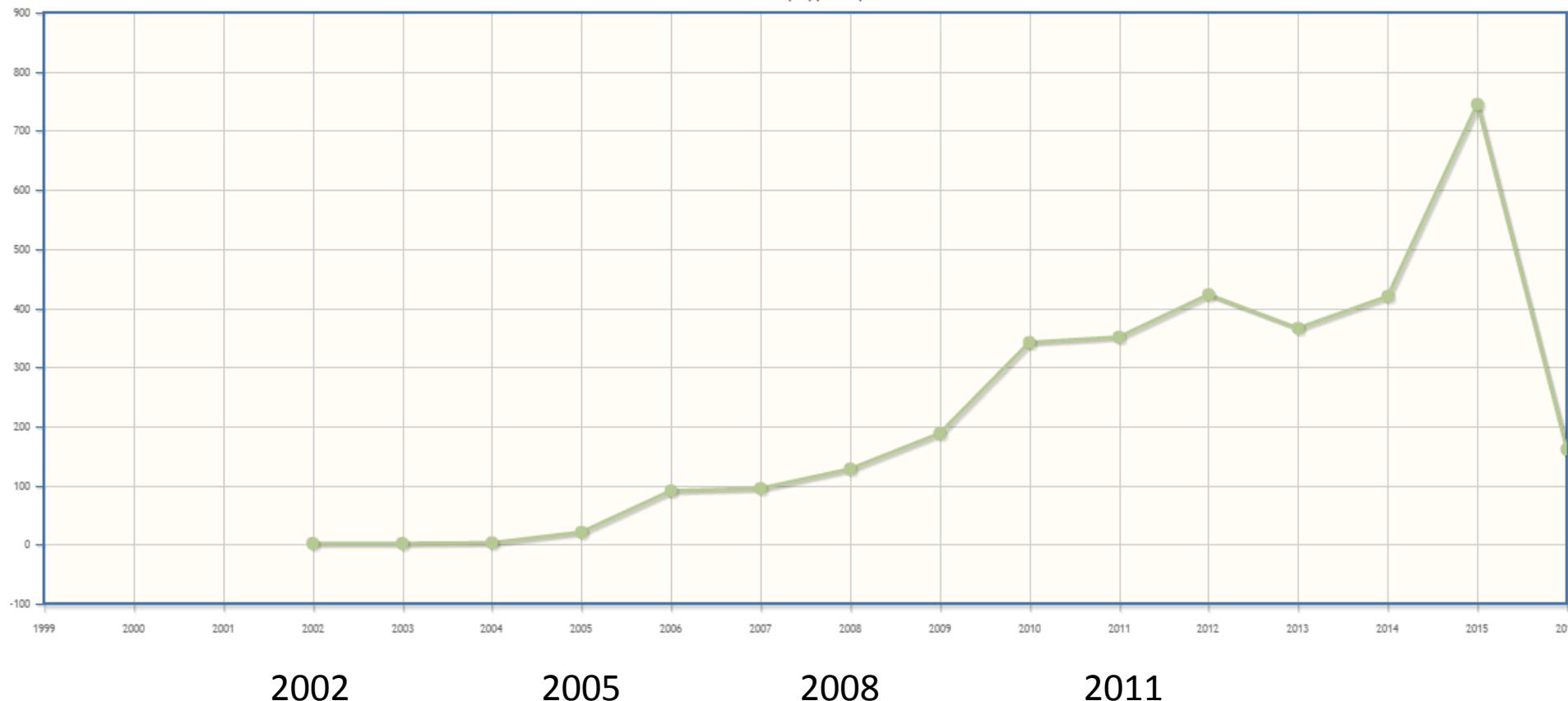
Compiled code



Main causes and exploitation methods

- Lack of input validation within methods (strcpy, gets, etc): stack based overflow (placing harmful code to the stack, ROP, JOP)
- Dynamic memory allocation problems (use after free, double free vulnerabilities) heap overflow (function pointer overwrite + heap spray)
- Exception handling errors (SEH overwrite)
- Others

Memory corruption vulnerabilities since 2002



What's the problem with this? (stack overflow)

```
#include <string.h>
void func1(char* ar1)
{
    char ar2[10];
    strcpy(ar2,ar1);
}
int main(int argc, char* argv[])
{
    func1(argv[1]);
}
```

What's the problem with this? (format string)

```
#include <string.h>
void func1(char* a, char* b)
{
    printf(a);
}

int main(int argc, char* argv[])
{
    func1(argv[1]);
}
```

What's the problem with this? (integer overflow)

```
if (channelp) {  
    /* set signal name (without SIG prefix) */  
    uint32_t namelen =  
        _libssh2_ntohu32(data + 9 + sizeof("exit-signal"));  
    channelp->exit_signal =  
        LIBSSH2_ALLOC(session, namelen + 1);  
    [...]  
    memcpy(channelp->exit_signal,  
           data + 13 + sizeof("exit_signal"), namelen);  
    channelp->exit_signal[namelen] = '\0';
```

What's the problem with this? (use after free)

```
char* ptr = (char*)malloc (SIZE);
if (err) {
    abrt = 1;
    free(ptr);
}
...
if (abrt) {
    logError("operation aborted before commit", ptr);
}
```

What's the problem with this? (double free)

```
char* ptr = (char*)malloc (SIZE);
```

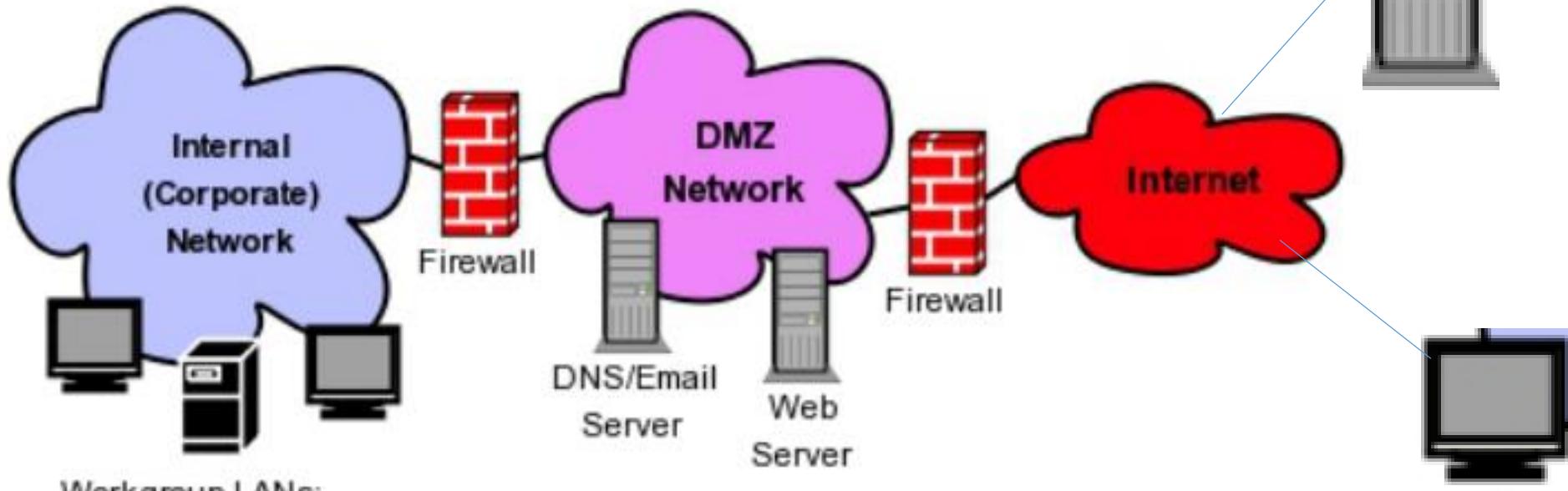
```
...
```

```
if (abrt) {  
    free(ptr);  
}  
  
...  
free(ptr);
```

Exploit dropper

Command & control

Typical network layout



Workgroup LANs:
Servers, workstations

Attacker's pc

Classic example of buffer overflow

...

Method1(a)

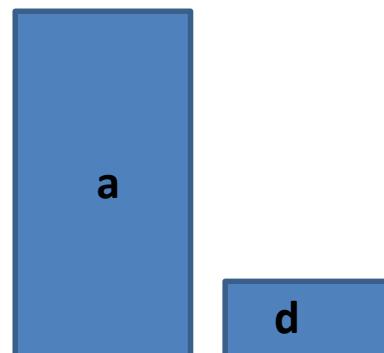
```
{  
d : fixed size array  
copy a to d  
}
```

Method2()

```
{  
Method1(a);  
}  
...
```

Code segment

```
00401280 | . E8 96610000 CALL <JMP.&CRTDLL.__GetMainArgs>  
00401292 | . B9 58804000 MOV ECX,OFFSET 00408058  
00401297 | . 8B11 MOV EDX,DWORD PTR DS:[ECX]  
00401299 | . 09D2 OR EDX,EDX  
0040129B | . 74 02 JZ SHORT 0040129F  
0040129D | . FFD1 CALL ECX  
0040129F > FF35 30A04000 PUSH DWORD PTR DS:[40A030]  
004012A5 | . FF35 2CA04000 PUSH DWORD PTR DS:[40A02C]  
004012AB | . FF35 28A04000 PUSH DWORD PTR DS:[40A028]  
004012B1 | . 8925 14A04000 MOV DWORD PTR DS:[40A014],ESP  
004012B7 | . E8 18000000 CALL 004012D4  
004012BC | . 83C4 18 ADD ESP,18  
004012BF | . 31C9 XOR ECX,ECX  
004012C1 | . 894D FC MOV DWORD PTR SS:[LOCAL.1],ECX  
004012C4 | . 50 PUSH EAX  
004012C5 | . E8 82610000 CALL <JMP.&CRTDLL.exit>  
004012CA | . C9 LEAVE  
004012CB | . C3 RETN
```



Stack

Method parameters

Return address

Saved frame pointer

Local variables

Method parameters

Return address

Saved frame pointer

Local variables

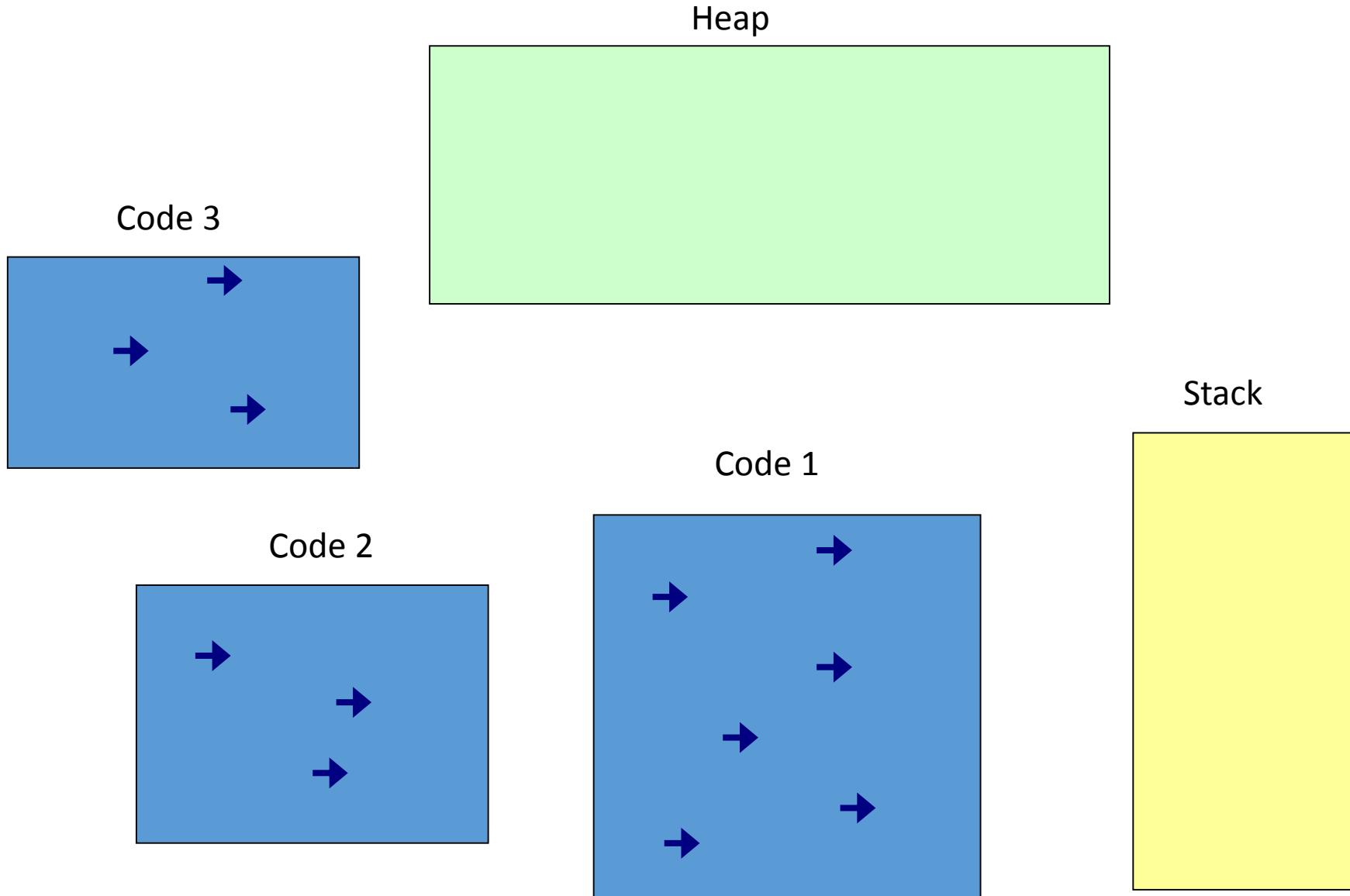
Method parameters

Return address

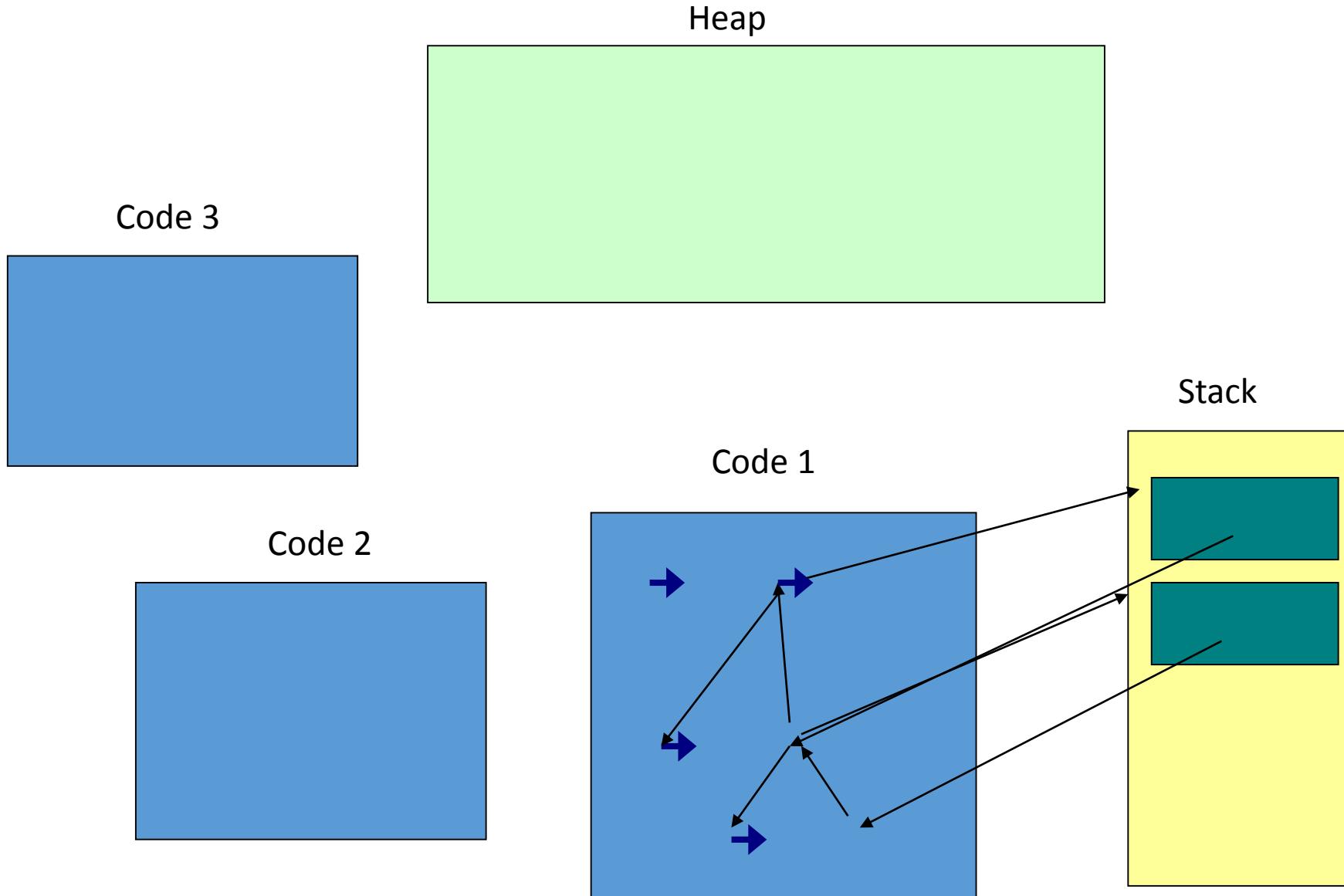
Saved frame pointer

Local variables

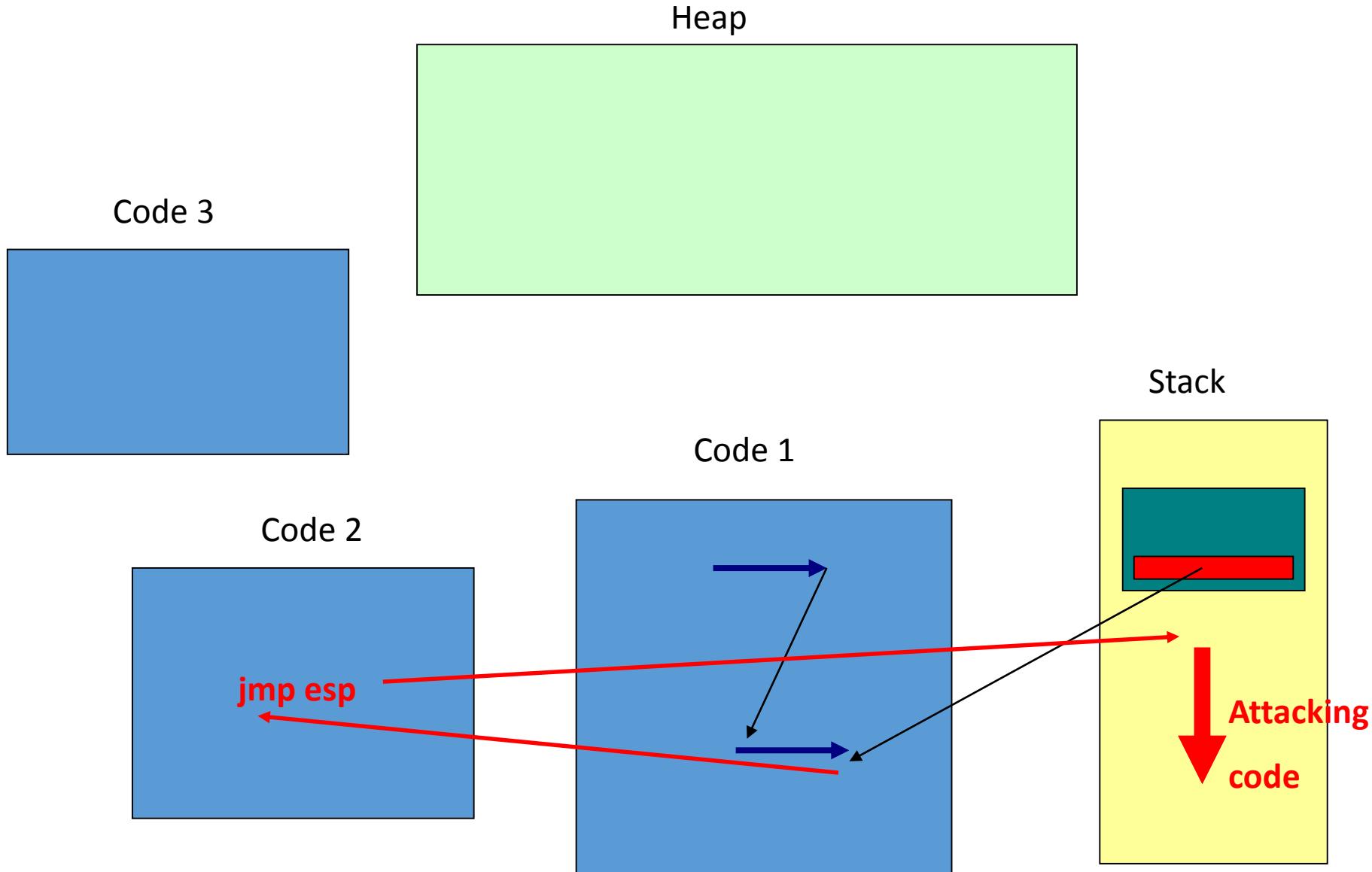
Normal operation



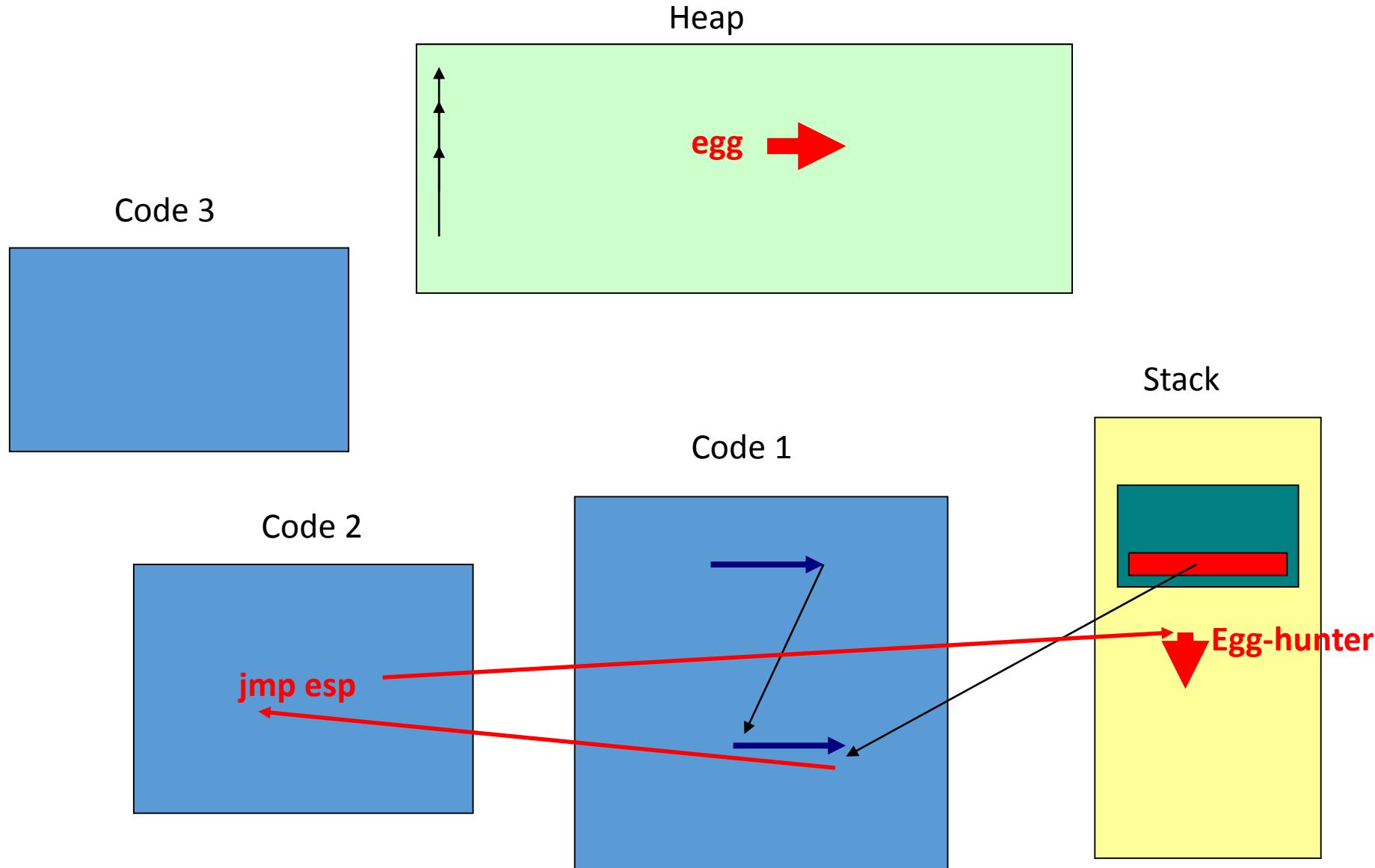
Normal operation



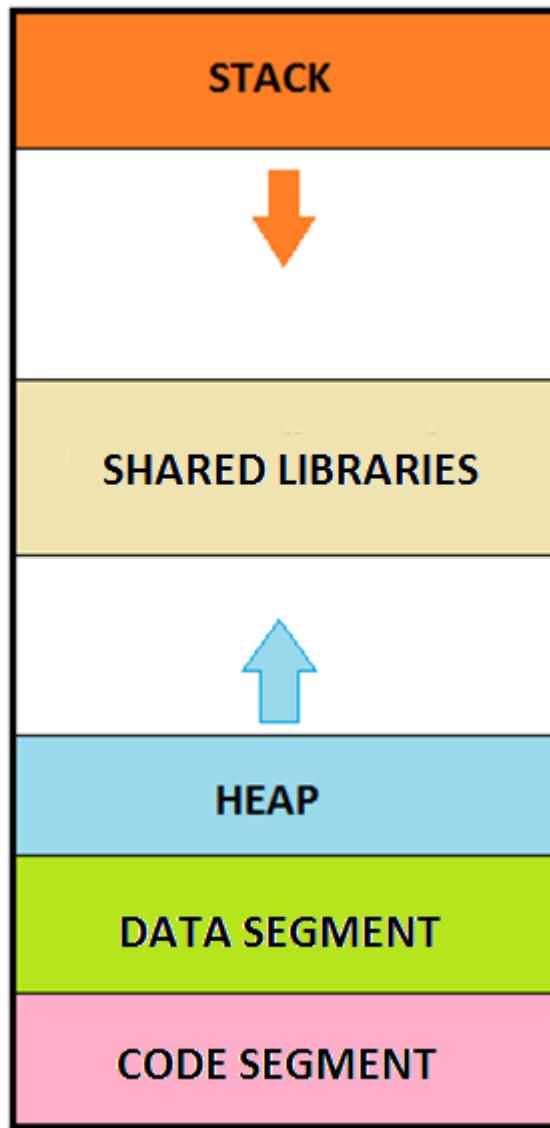
Normal operation



Egg-hunter



Data Execution Prevention



Data: read/write

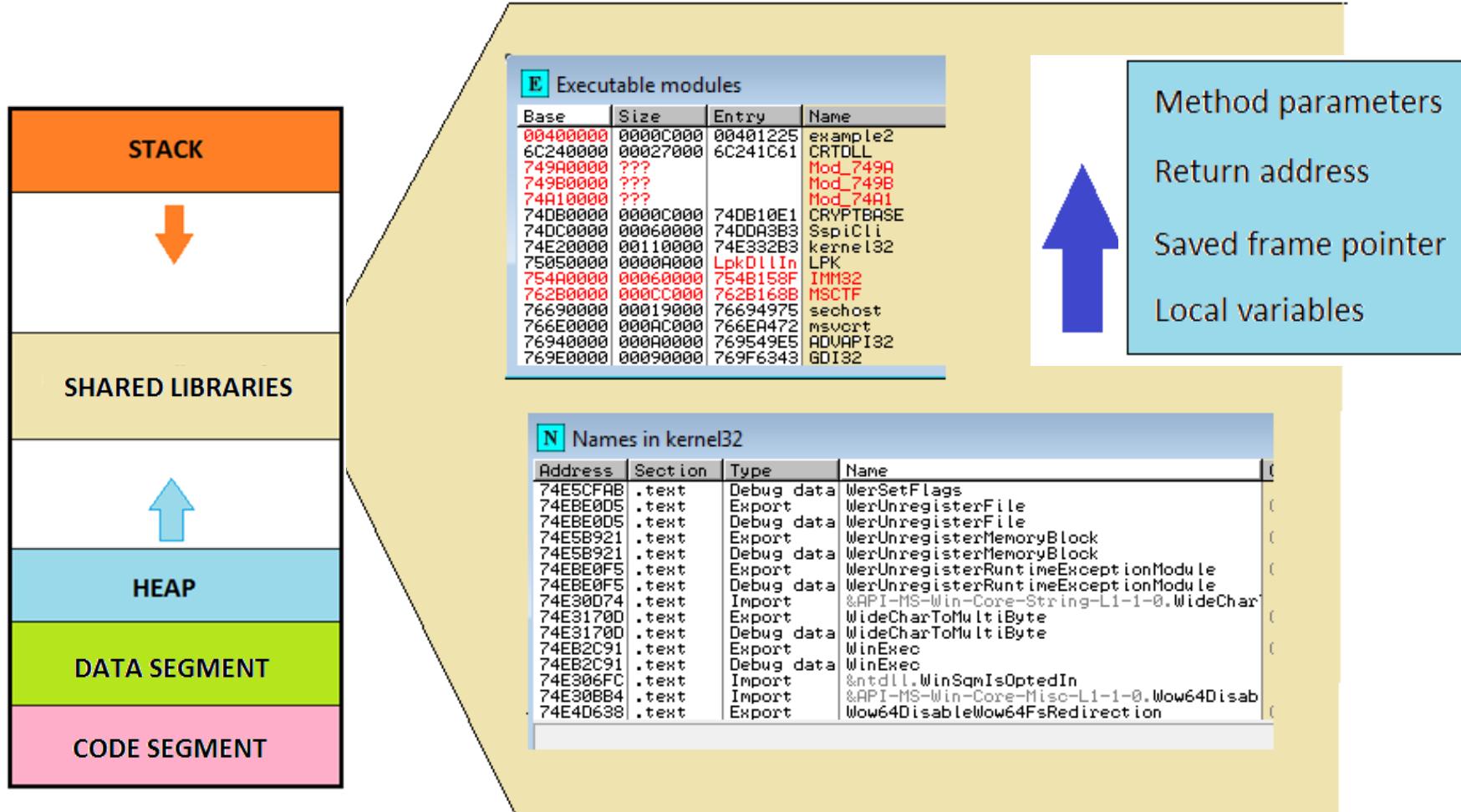
Code: read/execute

Data: read/write

Data: read/write

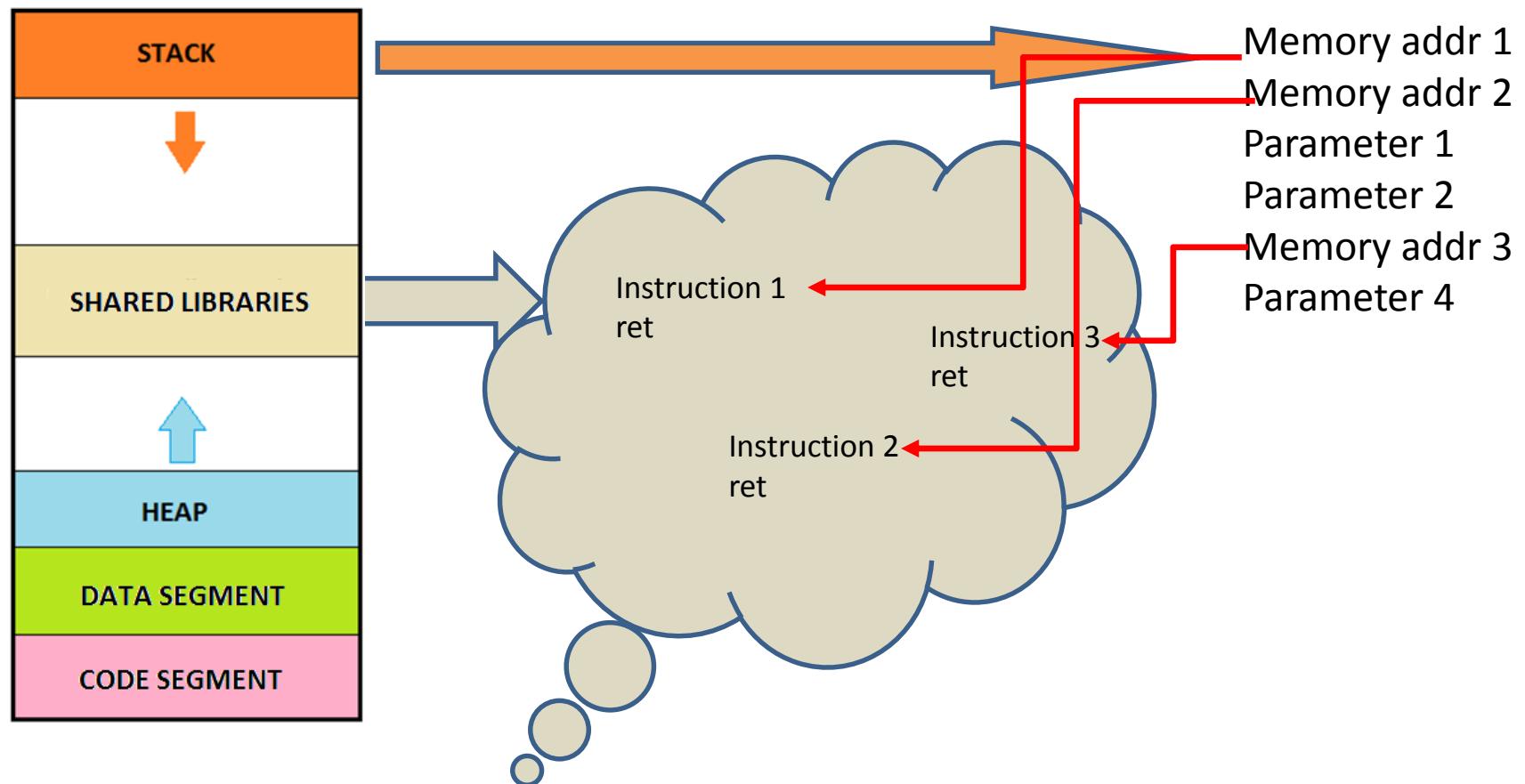
Code: read/execute

Avoiding memory execution protection (return to libc)

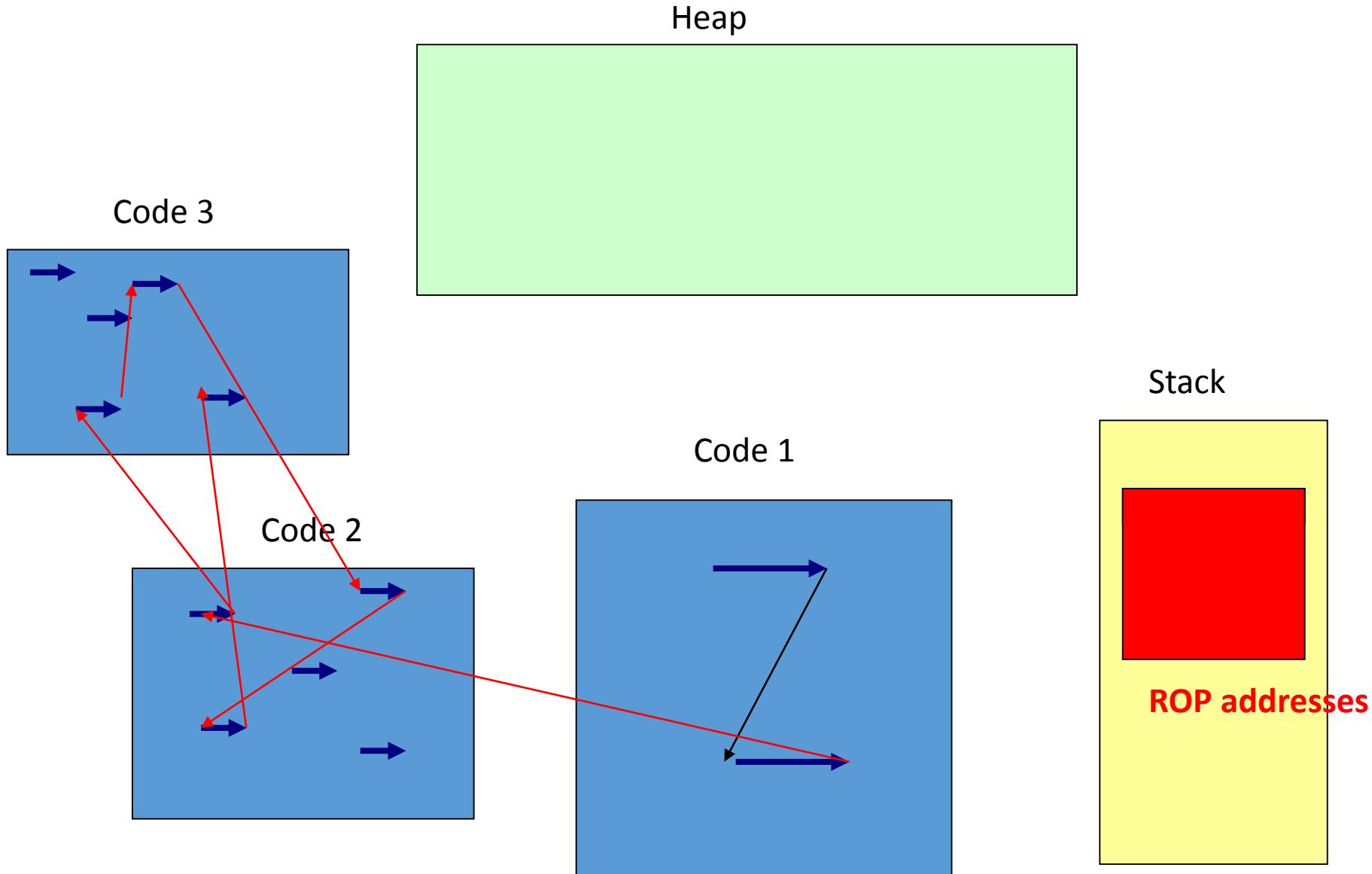


Avoiding DEP: Return oriented programming (ROP) Shacham, 2007

Executable code will not be placed on the stack only series of memory addresses and parameters



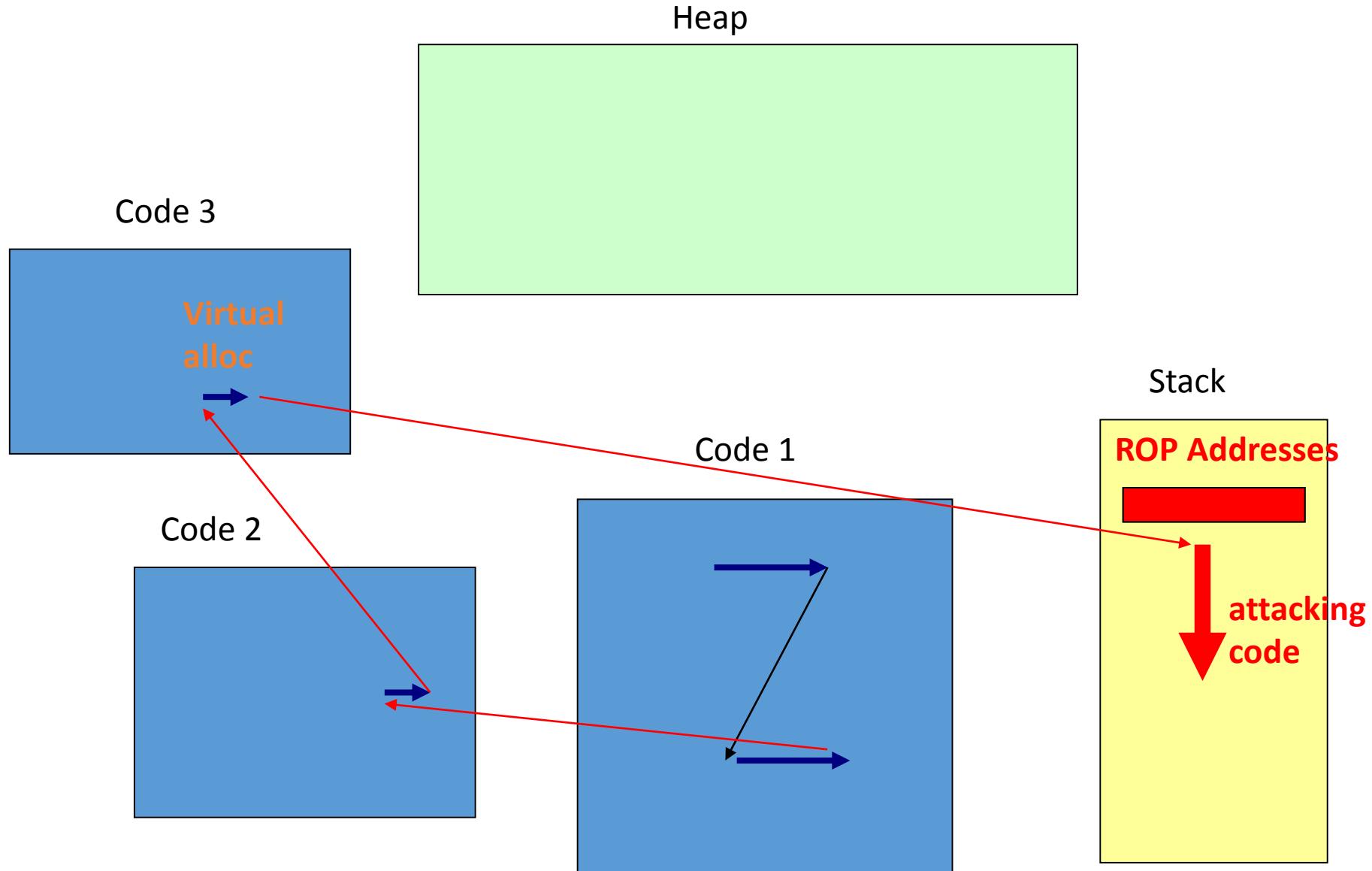
Return-Oriented Programming



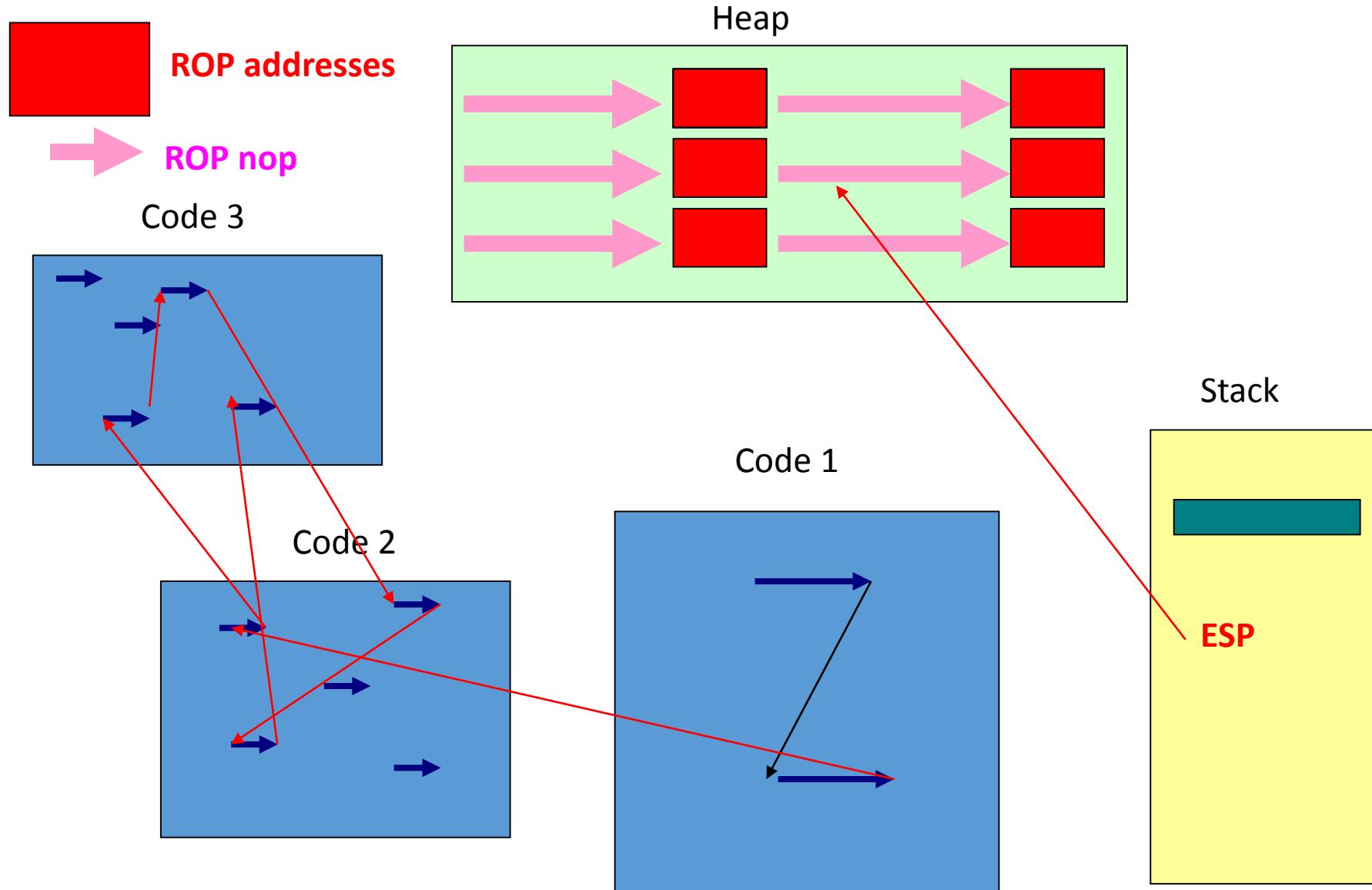
ROP – Turing completeness

- Instruction sequences
- Storing / loading variable
- If statement
- Loop execution
- Method call
- etc

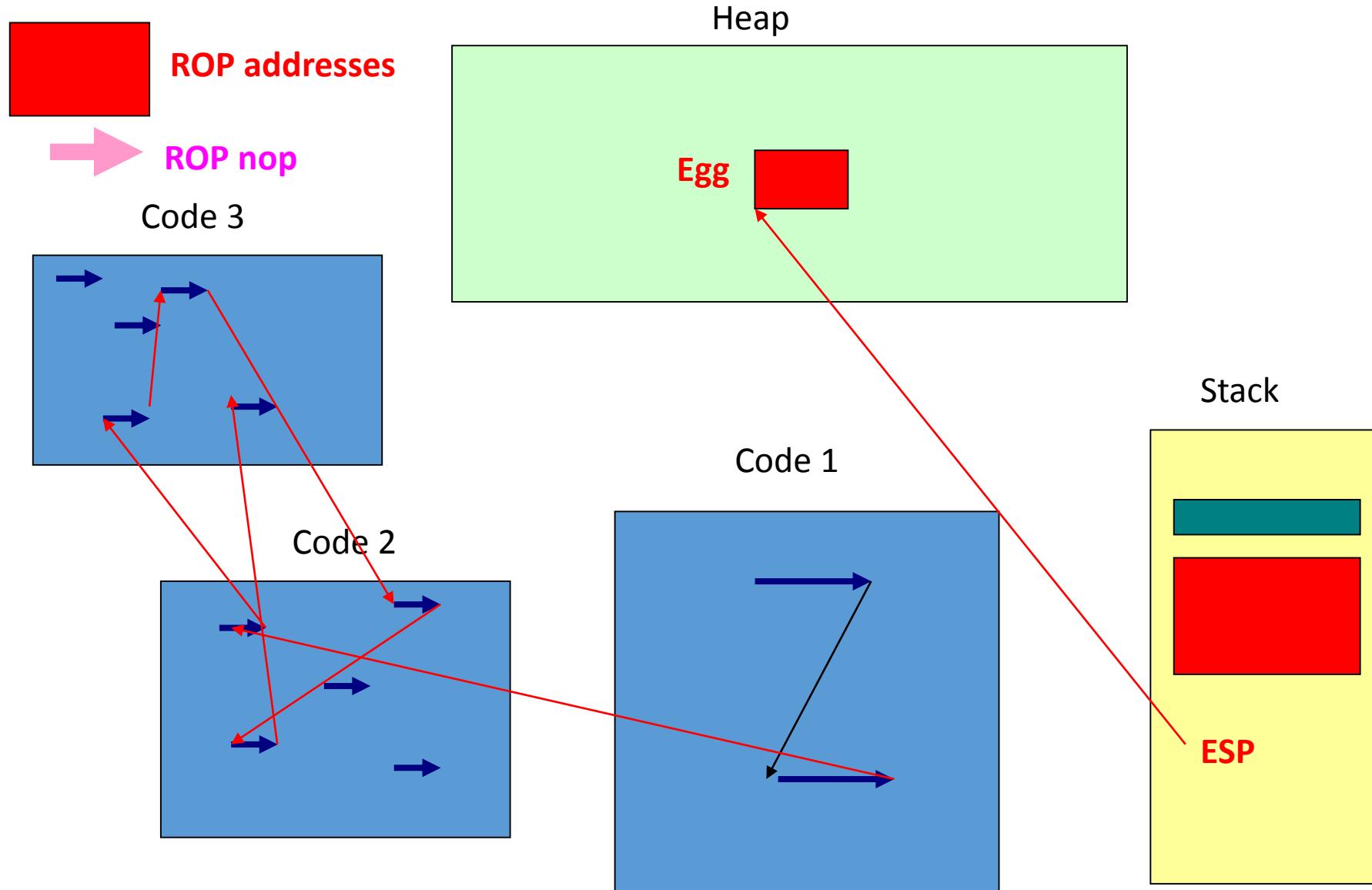
ROP + turn off DEP



ROP + Heap spray



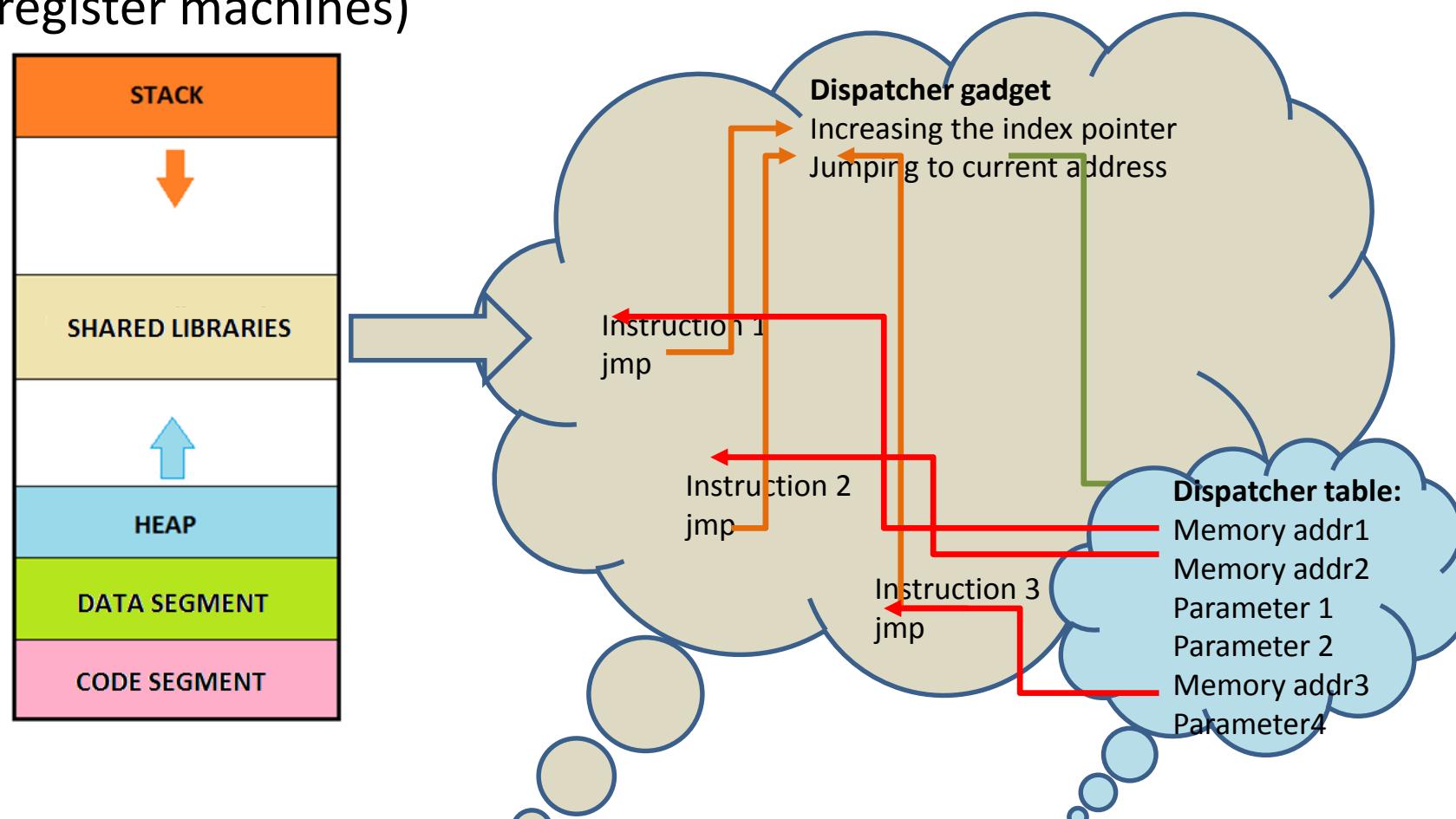
ROP + Egg-hunter



Jump oriented programming (JOP)

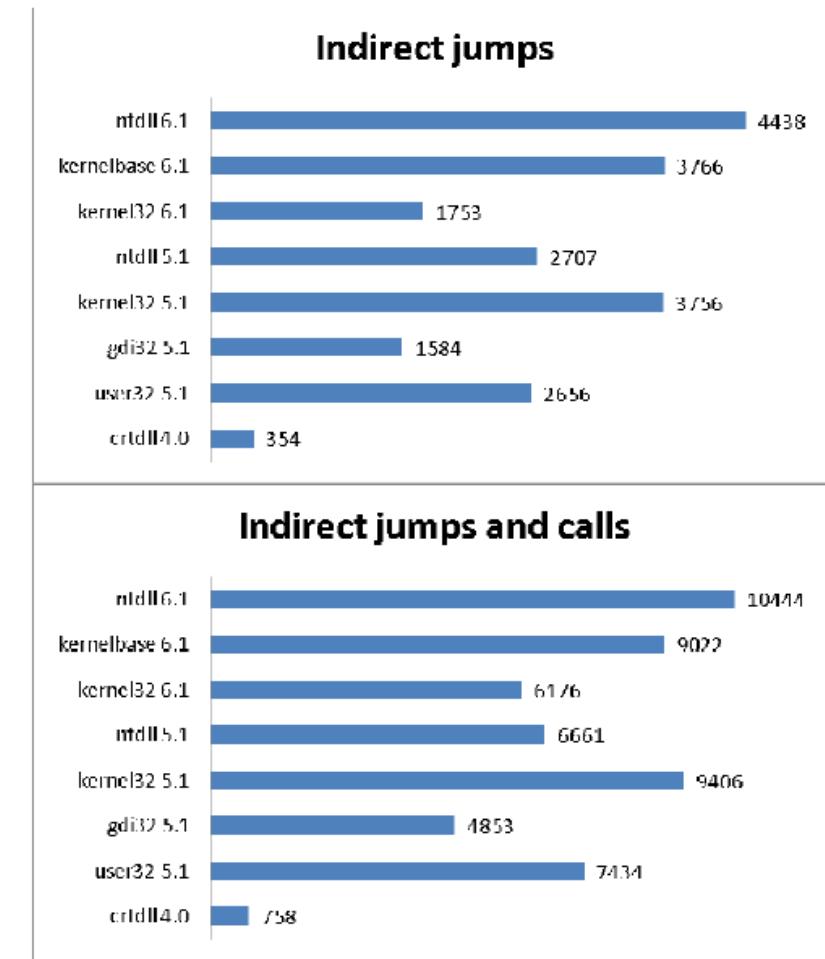
Bletsch, Jiang, Freeh 2011

- Attack execution without using stack (not sensible for stack cookie and returnless kernel, it can be used in the case of register machines)



Jump Oriented Programming – dispatcher gadgets in shared libraries (Erdődi, 2013)

File	Address	Opcode
crt.dll.dll 5.1.2600	73d3a066	add ebx,0x10 jmp dword ptr ds:[ebx]
crt.dll.dll 5.1.2600	73d3a0f2	add ebx,0x10 jmp dword ptr ds:[ebx]
user32.dll 5.1.2600	77d63ae9	add esi,edi jmp dword near [esi-0x75]
ntdll.dll 5.1.2600	7c939bbd	add ebx,0x10 jmp dword near [ebx]
ntdll.dll 5.1.2600	7c93c4db	sub edi,ebp call dword near [edi-0x18]
kernelbase. dll 6.2	75e6e815	sub esi,edi call dword near [esi+0x53]
ntdll.dll 6.2	77c94142	add ebx,0x10 jmp dword near [ebx]
ntdll.dll 6.2	77ca8c9	add ecx,edi jmp dword near [ecx+0x30]
ntdll.dll 6.2	77ca9dc0	add eax,edi call dword near [eax-0x18]
ntdll.dll 6.2	77cbcaca	add ebx,edi call dword near [ebx+0x5f]

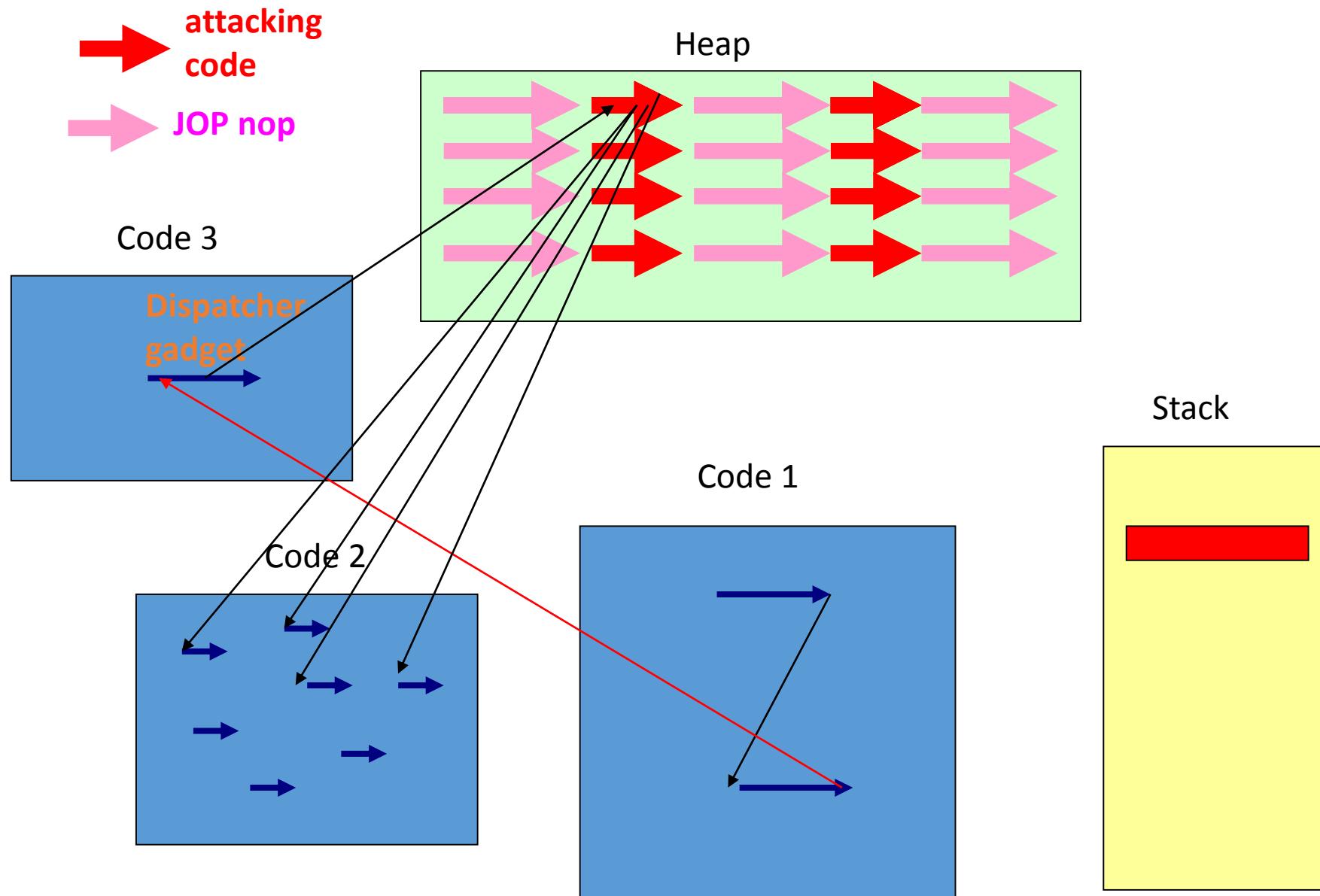


Jump Oriented Programming – WinExec example for Win32 X86

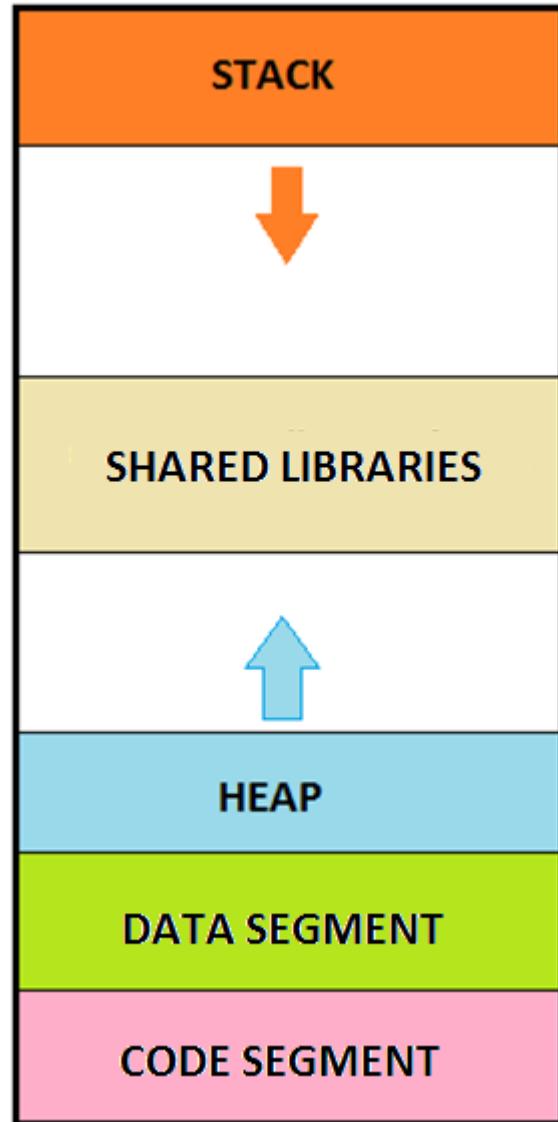
Address from the beginning of the dispatcher table	Value	Opcode	Function
0x00	77d65dda	pop eax std jmp ecx	sets eax to WinExec
0x10	77d5fa07	add esi,edi jmp ecx	sets esi to command string
0x20	77d482f6	xor edi,edi jmp ecx	zero edi
0x30	7c81ebb8	push edi jmp ecx	push zero on the stack
0x40	77d62d94	push esi std jmp ecx	push command string on the stack
0x50	7c9409ce	xchg esi,eax std jmp ecx	sets esi to WinExec

0x60	7c8306f0	mov edi,ebp jmp ecx	sets edi to dispatcher gadget
0x70	77f45ce1	call esi jmp edi	execute WinExec
0x80	77d482f6	xor edi,edi jmp ecx	zero edi
0x90	7c81ebb8	push edi jmp ecx	push zero on the stack
0xa0	77d65dda	pop eax std jmp ecx	sets eax to ExitProcess
0xb0	7c9409ce	xchg esi,eax std jmp ecx	sets esi to ExitProcess
0xc0	7c8306f0	mov edi,ebp jmp ecx	sets edi to dispatcher gadget
0xd0	77f45ce1	call esi jmp edi	execute ExitProcess

JOP + Heap spray



Address Space Layout Randomization (ASLR)



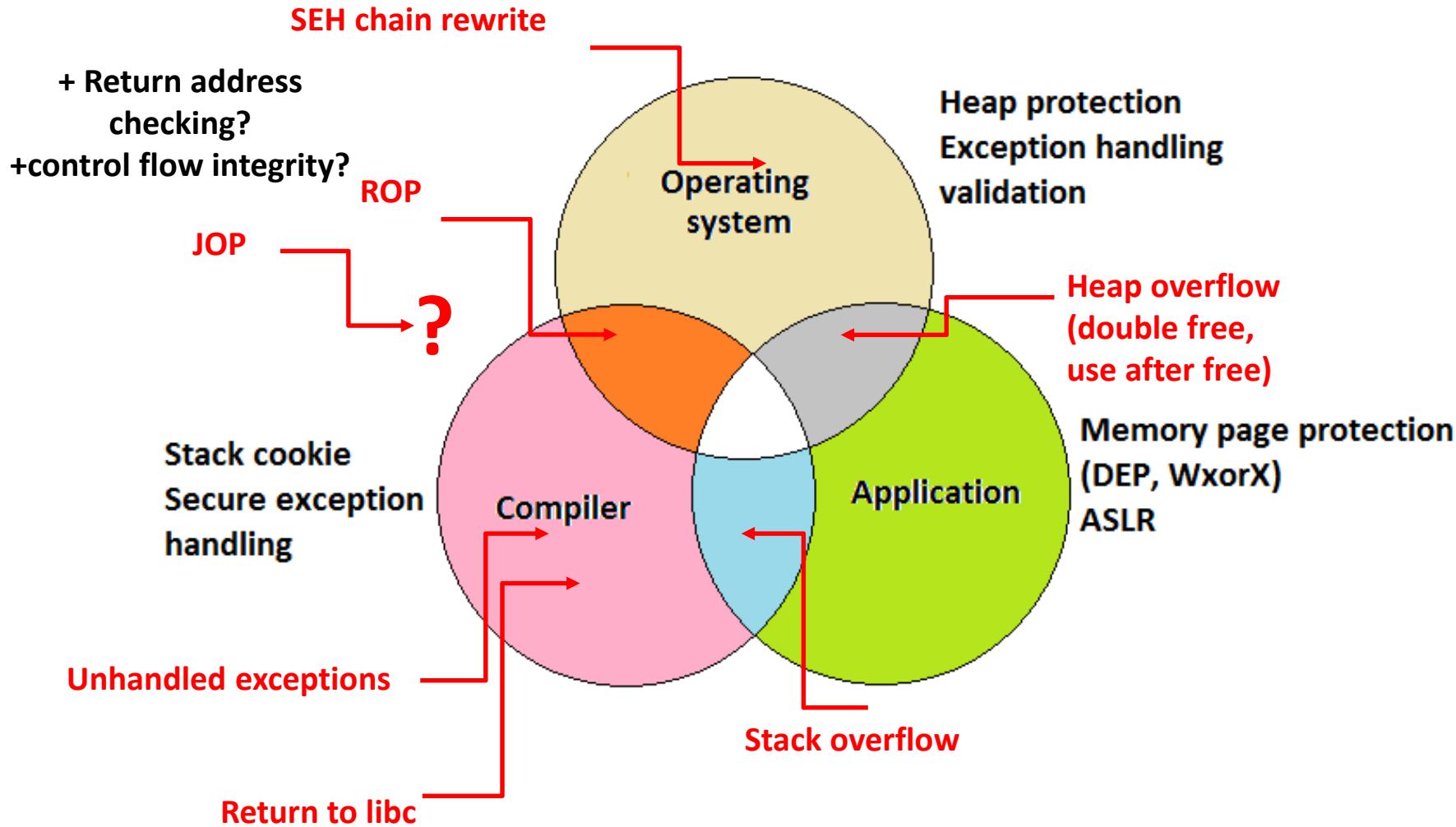
Bypassing ASLR

- Non Position Independent code segments
- Guessing the ASLR offset
- Information leakage
- JIT-ROP
- Blind ROP

Additional protections

- Windows Enhanced Mitigation Experience Toolkit (EMET)
- Execute no read (XnR)
- Returnless kernel?
- Return Address Checking
- Control Flow Integrity

Protection against memory corruption



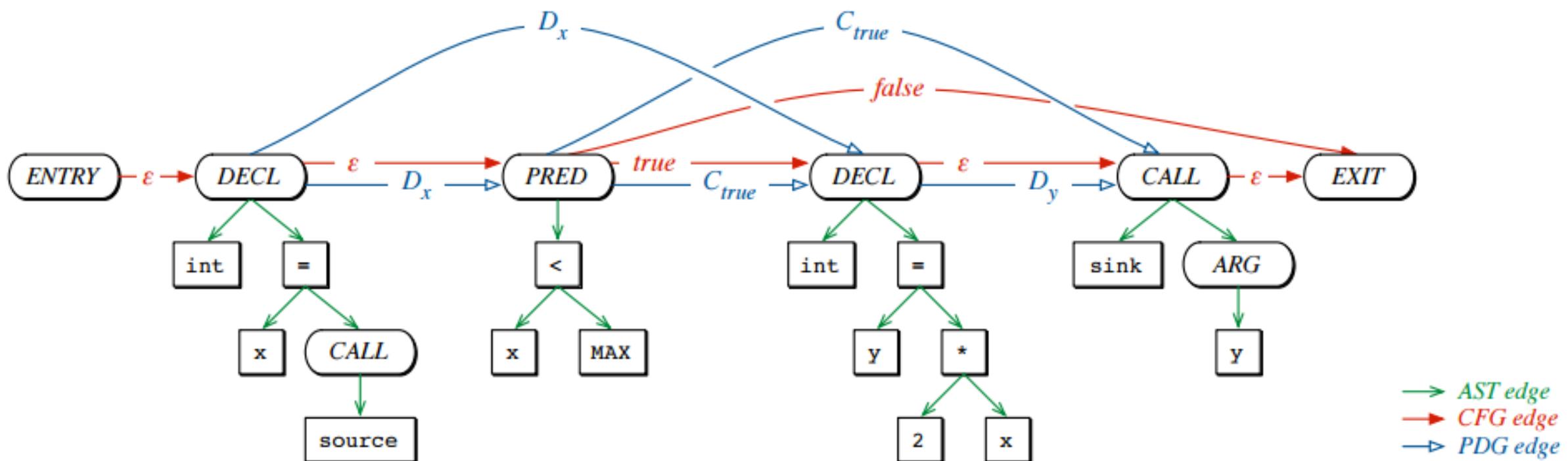
Vulnerability searching

- Static Analysis (source validators, Interactive Disassembler (IDA))
- Dynamic Analysis (Fuzzing)
- Finding vulnerability accidentally
- AV softwares by behaviour analysis (for already discovered non-public 0days)

Static code analyzers

- Unreachable codes
- Code duplicates
- Inappropriate memory management
- Lack of validation
- Etc.

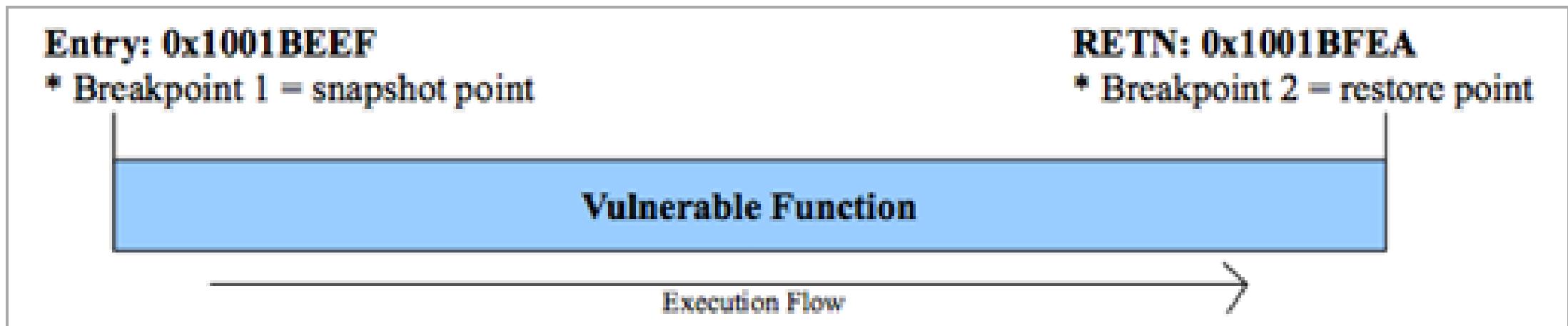
Code Property Graph (Yamaguchi et al, 2014)



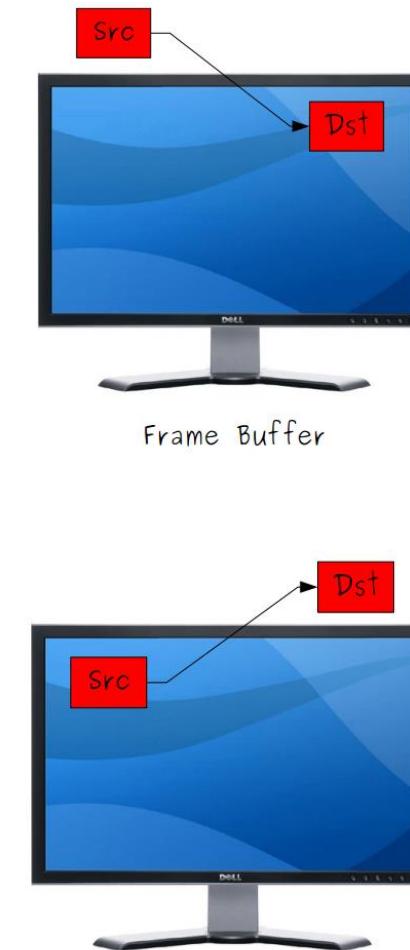
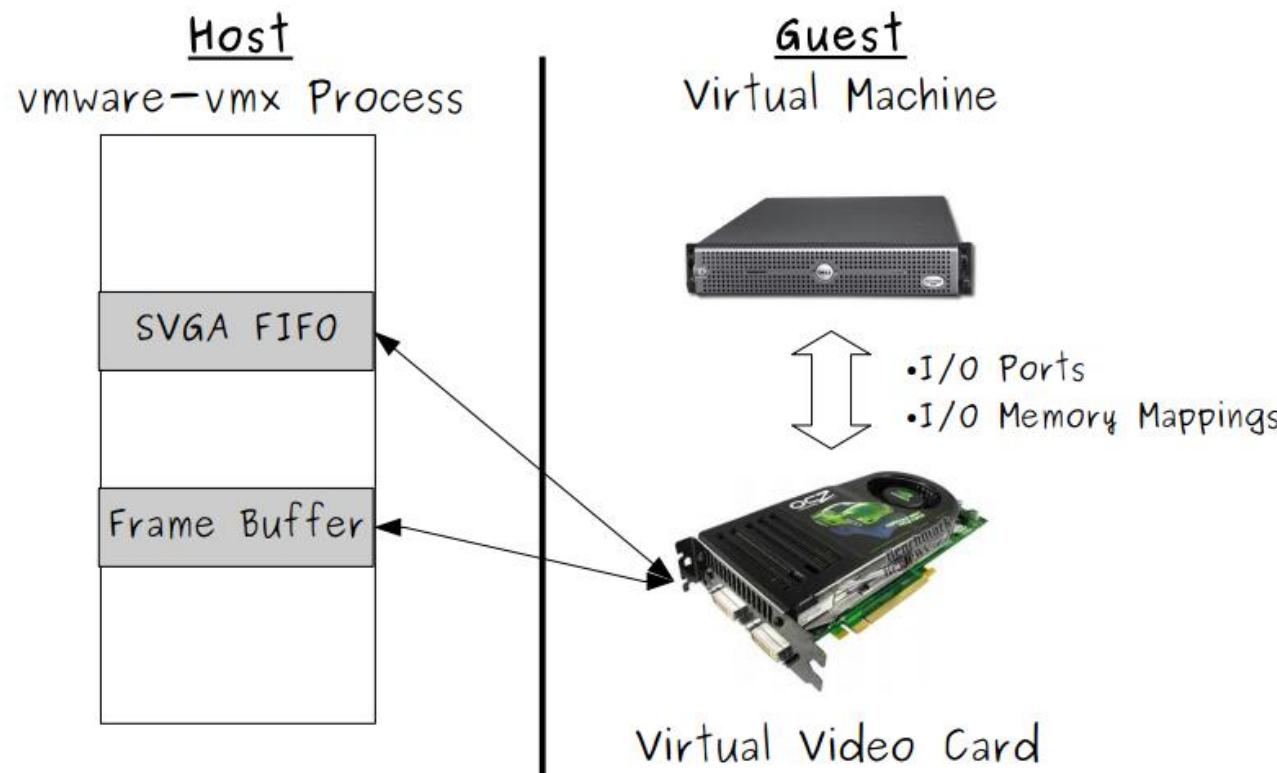
Input parameter / file format fuzzing

```
<?xml version="1.0" encoding="utf-8"?>
<Peach xmlns="http://peachfuzzer.com/2012/Peach" xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
       xsi:schemaLocation="http://peachfuzzer.com/2012/Peach ../peach.xsd">
    <Include ns="Http" src="file:http_base.xml" />
    <StateModel name="RandomFuzzing" initialState="initialRandomFuzzing">
        <State name="initialRandomFuzzing">
            <Action type="output">
                <DataModel name="Request" ref="Http:Request" />
                <Data>
                    ...
                    ...
                </Data>
            </Action>
            <Action type="input">
                <DataModel name="Response" ref="Http:Response" />
            </Action>
            ...
            ...
        </State>
    </StateModel>
    <Test name="Default">
        <Agent name="Ping-Agent">
            <Monitor class="Ping">
                ...
            </Monitor>
        </Agent>
        <Strategy class="RandomDeterministic" />
        <StateModel ref="RandomDeterministicFuzzing" />
        <Publisher class="TcpClient">
            ...
        </Publisher>
        <Logger class="File">
            <Param name="Path" value="logs" />
        </Logger>
    </Test>
</Peach>
```

In memory fuzzing



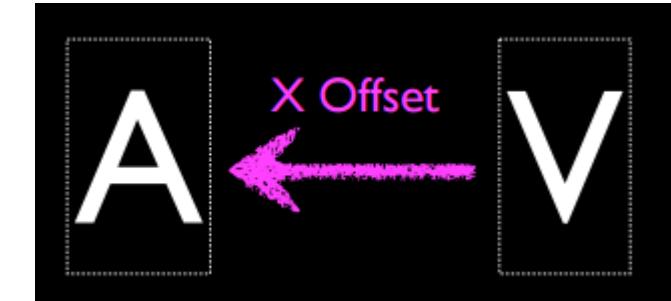
Example memory corruption: Cloudburst (Kortchinsky, 2009)



Example memory corruptions (TrueType Font Engine Vulnerability)

Font Program

b8 = PUSHW	00 04 00 00 45 42 44 54	EBDT
7fc0 = 32704	00 00 00 28 45 42 4c 43	K.C...T...(EBLC
b8 = PUSHW	00 00 01 78 45 42 53 43	.M2.... ...xEBSC
000000 01c0 = 448	00 00 00 94 4f 53 2f 32OS/2
000000 63 = MUL	00 00 00 56 63 6d 61 70\$...Vcmap
000000 b8 = PUSHW	00 00 00 34 63 76 74 20	a.W.....4cvt
000000 3a40 = 14912	00 00 00 02 66 70 67 6dfpgm
000000 60 = ADD	00 03 b8 9b 67 6c 79 66glyf
	00 00 00 bc 68 65 61 64	.iK.....head
	00 00 00 36 68 68 65 61	..(.....6hhea
	00 00 00 24 68 6d 74 78\$hmtx
	00 00 00 0e 6c 6f 63 61loca
	00 00 00 0e 6d 61 78 70	.^.....maxp
	00 00 00 20 6e 61 6d 65	...#.....name
	00 00 01 7c 70 6f 73 74 post
	00 00 00 35 70 72 65 70	..>i.....5prep
000000100 8b 9d ff 81 00 03 ba c0	00 00 00 0d b8 7f c0 b8
000000110 01 c0 63 b8 3a 40 60 b8	00 0c 60 1c 00 00 00 00	..c.:@`....
000000120 00 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00
etc...		



Thank you!